

# Document made available under the Patent Cooperation Treaty (PCT)

International application number: PCT/EP04/053447

International filing date: 14 December 2004 (14.12.2004)

Document type: Certified copy of priority document

Document details: Country/Office: EP  
Number: 04100811.1  
Filing date: 01 March 2004 (01.03.2004)

Date of receipt at the International Bureau: 09 March 2005 (09.03.2005)

Remark: Priority document submitted or transmitted to the International Bureau in compliance with Rule 17.1(a) or (b)



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**Patentanmeldung Nr.    Patent application No.    Demande de brevet n°**

**04100811.1**

Der Präsident des Europäischen Patentamts;  
Im Auftrag

For the President of the European Patent Office

Le Président de l'Office européen des brevets  
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**R C van Dijk**





Anmeldung Nr:  
Application no.: 04100811.1  
Demande no:

Anmeldetag:  
Date of filing: 01.03.04  
Date de dépôt:

Anmelder/Applicant(s)/Demandeur(s):

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Bezeichnung der Erfindung/Title of the invention/Titre de l'invention:  
(Falls die Bezeichnung der Erfindung nicht angegeben ist, siehe Beschreibung.  
If no title is shown please refer to the description.  
Si aucun titre n'est indiqué se référer à la description.)

Method and device for transcoding N-bit words into M-bit words with M N

In Anspruch genommene Priorität(en) / Priority(ies) claimed /Priorité(s)  
revendiquée(s)  
Staat/Tag/Aktenzeichen/State/Date/File no./Pays/Date/Numéro de dépôt:

EP/18.12.03/EP 03293224

Internationale Patentklassifikation/International Patent Classification/  
Classification internationale des brevets:

G09G/

Am Anmeldetag benannte Vertragsstaaten/Contracting states designated at date of  
filing/Etats contractants désignées lors du dépôt:

AT BE BG CH CY CZ DE DK EE ES FI FR GB GR HU IE IT LU MC NL  
PL PT RO SE SI SK TR LI



## METHOD AND DEVICE FOR TRANSCODING N-BIT WORDS INTO M-BIT WORDS WITH $M < N$

5 The present invention concerns a method for transcoding a N bits word into a M bits word with  $M < N$ . The invention is applicable in various fields and more particularly in the display field.

### Background

10 The invention will be described in the field of the plasma display device. The transcoding method is used after a degamma function in a display receiver. Figure 1 is a standard implementation circuit for encoding video data into subfield data. This circuit comprises a first memory 100 comprising 1024 x 12 bits for handling 10 bits video data  $YI[9:0]$ . A first look up table is stored in the memory 100 and used for remapping these data as described in EP  
15 patent application No. 1 353 314 for implementing the metacode concept or in EP patent application No 1 256 954 for implementing a Gravity Center Coding. At the output of the memory 100, 12 bits video signal  $YA[11:0]$  is obtained. The available 12-bits correspond to 8-bits integer resolution and 4-bits fractional resolution. Then, the 12-bits of video signal  $YA [11-0]$  are  
20 forwarded to a dithering circuit 110. In this circuit 110, the 4-bits of fractional resolution are added with the 4-bits of dithering and then truncated.

The video signal  $YB[7,0]$  from the circuit 110 is then forwarded to a second memory 120 comprising 256 x 16 bits. A second look-up table is stored in the memory 120 and is used for transforming the data  $YB[7:0]$  into  
25 subfield data  $SF[15:0]$ .

Until now, a degamma function is implemented in the look-up table of the memory 100. Otherwise, in order not to lose significant bits, much more bits (typically 16 bits) would be required after the degamma function than the video input has. So, if the degamma function is made  
30 before the memory 100, this memory must have more inputs (16 bits), which will drastically increase the size of this memory. To go from 10 to 16 bits, the size of the memory is multiplied by 64.

And, on another hand, if a user wants to change the value of the degamma, to use a different degamma for each color or to use a specific degamma function, a specific look-up table should be generated for each case (if the degamma function can be known) and all of them should be  
 5 stored.

For these reasons, it is necessary to separate the degamma function from the look-up table of the memory 100 and to create a transcoding circuit for transcoding 16 bits data (coming for example from a degamma block) into 10 bits data.

10 An object of the invention is to propose a method for transcoding a 16 bits word into a 10 bits word and, more generally, a method for transcoding a N bits word into a M bits word,  $N > M$ .

#### Invention

15 The invention concerns a method for transcoding a N bits word into a M bits word,  $N > M$ , characterized in that it comprises the following steps :

- breaking down the N bits word into an exponent part and a mantissa part having each a size which varies according to the value of said  
 20 N bits word, the size of the mantissa part increasing with the value of said N bits word, and
- encoding the exponent part of the N bits word into a variable number of bits A and removing, if need be, least significant bits of the mantissa part in order to obtain a mantissa with a variable number of bits B,  
 25 with  $A+B=M$ .

The invention concerns also a device for transcoding a N bits word into a M bits word,  $N > M$ , characterized it comprises

- means for breaking down the N bits word into an exponent part  
 30 and a mantissa part having each a size which varies according to the value of said N bits word, the size of the mantissa part increasing with the value of said N bits word, and
- means for encoding the exponent part of the N bits word into a variable number of bits A and removing, if need be, least significant bits of

the mantissa part in order to obtain a mantissa with a variable number of bits B, with  $A+B=M$ .

5 It concerns also a Plasma display panel comprising a degamma means for applying a degamma function to video input data and delivering N bits data and a mapping memory for remapping M bits data, with  $N>M$  characterized in that it comprises a transcoding device as above.

#### drawings

10 Exemplary embodiments of the invention are illustrated in the drawings and are explained in more detail in the following description, the drawings showing in :

Fig.1 a block diagram of a standard means for encoding video data into subfield data; and

15 Fig.2 a block diagram of a means for encoding video data into subfield data with an transcoding block according to the invention;

#### exemplary embodiments

20 The method is described for transcoding a 16 bits word into a 10 bits word. As this transcoding method is used for the output of a degamma function, the input values of this transcoding method represent linear video information. So since the eye is less sensitive in the bright levels than in the darks levels, the Least Significant Bits of a 16 bits word representing a high  
25 value are not important (and may be even not significant). This means that high values do not require as much precision as low values. On the other hand, the Most Significant Bits of a 16 bits word representing a low level are not necessary since they are all equal zero. This means that low values do not require the 16 bits.

30 For example, for a 16 bits word, values inferior to 256 do not need more than 8 bits to be encoded without loss of accuracy and values superior to 1024 do not need to have a 1-bit precision. Consequently, some bits of the word could be saved.

35 According to the invention; each value is defined by a mantissa and an exponent. So, the 10 bit word comprises a mantissa part and an exponent



part. The idea of the invention is to use a reduced number of bits for the mantissa.

Basically, it is used 8 bits for the mantissa part, except for the smallest values for which less bits are used. In fact, it is equivalent to use for the mantissa part more bits (up to 16) for the small values and less for the biggest ones (down to 8).

An example of such a conversion is given below. In this example, the 65536 16-bits input values are converted into 1024 output values (10 bits).

The 16-bits values can be classified as following :

	1XXXXXXXXXXXXXXXXX	32768<value<65535
	01XXXXXXXXXXXXXXXX	16384<value<32767
	001XXXXXXXXXXXXXX	8162<value<16383
15	0001XXXXXXXXXXXX	4096<value<8161
	00001XXXXXXXXXXX	2048<value<4095
	000001XXXXXXXXXX	1024<value<2047
	0000001XXXXXXXXX	512<value<1023
	00000001XXXXXXXX	256<value<511
20	000000001XXXXXX	128<value<255
	0000000001XXXXXX	64<value<127
	00000000001XXXXX	32<value<63
	00000000000XXXXX	0<value<31

According to the invention, the small x (symbolizing indifferently zeros or ones) on the right part are lost bits and the zeros and ones on the left side are information which are preserved by an encoding (they are defined by the exponent). The big x (symbolizing indifferently zeros or ones) are bits which are preserved as they are in the 16-bit word.

30

The transcoding operation can be illustrated by the two following table. The left table corresponds to the 16 bits words and the right table corresponds to the 10 bits words.

	1XXXXXXXXXXXXXXXXX	→	111XXXXXXXX
	01XXXXXXXXXXXXXXXX	→	110XXXXXXXX
	001XXXXXXXXXXXXXXXX	→	101XXXXXXXX
	0001XXXXXXXXXXXXXX	→	100XXXXXXXX
5	00001XXXXXXXXXXXXX	→	011XXXXXXXX
	000001XXXXXXXXXXX	→	010XXXXXXXX
	0000001XXXXXXXXXX	→	0011XXXXXXX
	00000001XXXXXXXXX	→	0010XXXXXXX
	000000001XXXXXXX	→	00011XXXXXX
10	0000000001XXXXXX	→	00010XXXXXX
	00000000001XXXXX	→	00001XXXXXX
	00000000000XXXXX	→	00000XXXXXX

- 15 In this example,
- the numbers between 32768 and 65535 are encoded with 7 original bits (32768 input values to 128 output values); the mantissa comprises 7 bits and the exponent comprises 3 bits.
  - the numbers between 16384 and 32767 are encoded with 7 original bits
  - 20 (16384 input values to 128 output values); the mantissa comprises 7 bits and the exponent comprises 3 bits.
  - the numbers between 8192 and 16383 are encoded with 7 bits (8192 input values to 128 output values); the mantissa comprises 7 bits and the exponent comprises 3 bits.
  - 25 - the numbers between 4096 and 8191 are encoded with 7 original bits (4096 input values to 128 output values) ; the mantissa comprises 7 bits and the exponent comprises 3 bits.
  - the numbers between 2048 and 4095 are encoded with 7 original bits (2048 input values to 128 output values) ; the mantissa comprises 7 bits and
  - 30 the exponent comprises 3 bits.
  - the numbers between 1024 and 2047 are encoded with 7 original bits (1024 input values to 128 output values); the mantissa comprises 7 bits and the exponent comprises 3 bits.
  - the numbers between 512 and 1023 are encoded with 6 original bits (512
  - 35 input values to 64 output values) ; the mantissa comprises 6 bits and the exponent comprises 4 bits.
  - the numbers between 256 and 511 are encoded with 6 original bits (256 input values to 64 output values) ; the mantissa comprises 6 bits and the exponent comprises 4 bits.

- the numbers between 128 and 255 are encoded with 5 original bits (128 input values to 32 output values) ; the mantissa comprises 5 bits and the exponent comprises 5 bits.
- the numbers between 64 and 127 are encoded with 5 original bits (64 input values to 32 output values) ; the mantissa comprises 5 bits and the exponent comprises 5 bits.
- the numbers between 32 and 63 are encoded with 5 original bits (32 input values to 32 output values) ; the mantissa comprises 5 bits and the exponent comprises 5 bits.
- the numbers between 0 and 31 are encoded with 5 original bits (32 input values into 32 output values); the mantissa comprises 5 bits and the exponent comprises 5 bits.

This encoding is done with some logic elements as indicated below. In[15:0] is the 16 bits word to be encoded and out [9:0] is the 10 bit word.

```

    if( in[15] == 1 )
        out[9:7] = 7
        out[6:0] = in[14:8]
    else if( in[14] == 1 )
        out[9:7] = 6
        out[6:0] = in[13:7]
    else if( in[13] == 1 )
        out[9:7] = 5
        out[6:0] = in[12:6]
    else if( in[12] == 1 )
        out[9:7] = 4
        out[6:0] = in[11:5]
    else if( in[11] == 1 )
        out[9:7] = 3
        out[6:0] = in[10:4]
    else if( in[10] == 1 )
        out[9:7] = 2
        out[6:0] = in[9:3]
    else if( in[9] == 1 )
        out[9:6] = 3
        out[5:0] = in[8:3]

```

```

else if( in[8] == 1 )
    out[9:6] = 2
    out[5:0] = in[7:2]
else if( in[7] == 1 )
5    out[9:5] = 3
    out[4:0] = in[6:2]
else if( in[6] == 1 )
    out[9:5] = 2
    out[4:0] = in[5:1]
10    else if( in[5] == 1 )
        out[9:5] = 1
        out[4:0] = in[4:0]
    else
15        out[9:5] = 0
        out[4:0] = in[4:0]

```

This transcoding method is implemented by a transcoding block 90  
 inserted between a degamma block 80 delivering 16 bits data and the  
 memory 100 as illustrated by figure 2. Of course, in this circuit, the look-up  
 20 table of the memory 100 does not include the degamma function but should  
 include a reverse transcoding of the block 90.



## CLAIMS

1. Method for transcoding a N bits word into a M bits word,  $N > M$ , characterized in that it comprises the following steps :

- 5                   - breaking down the N bits word into an exponent part and a mantissa part having each a size which varies according to the value of said N bits word, the size of the mantissa part increasing with the value of said N bits word, and
- 10                  - encoding the exponent part of the N bits word into a variable number of bits A and removing, if need be, least significant bits of the mantissa part in order to obtain a mantissa with a variable number of bits B, with  $A+B=M$ .

2. Device for transcoding a N bits word into a M bits word,  $N > M$ , characterized it comprises

- 15                   - means for breaking down the N bits word into an exponent part and a mantissa part having each a size which varies according to the value of said N bits word, the size of the mantissa part increasing with the value of said N bits word, and
- 20                   - means for encoding the exponent part of the N bits word into a variable number of bits A and removing, if need be, least significant bits of the mantissa part in order to obtain a mantissa with a variable number of bits B, with  $A+B=M$ .

25                  3. Plasma display panel comprising a degamma means (80) for applying a degamma function to video input data ( $Y[9:0]$ ) and delivering N bits data and a mapping memory (100) for remapping M bits data, with  $N>M$  characterized in that it comprises a transcoding device according to claim 2.

30



ABSTRACT**METHOD AND DEVICE FOR TRANSCODING N-BIT WORDS  
INTO M-BIT WORDS WITH  $M < N$** 

5

The present invention concerns a method for transcoding a N bits word into a M bits word with  $M < N$ . The invention is applicable in various fields and more particularly in the display field. The method comprises the following steps :

10

- breaking down the N bits word into an exponent part and a mantissa part having each a size which varies according to the value of said N bits word, the size of the mantissa part increasing with the value of said N bits word, and

15

- encoding the exponent part of the N bits word into a variable number of bits A and removing, if need be, least significant bits of the mantissa part in order to obtain a mantissa with a variable number of bits B, with  $A+B=M$ .





1/1

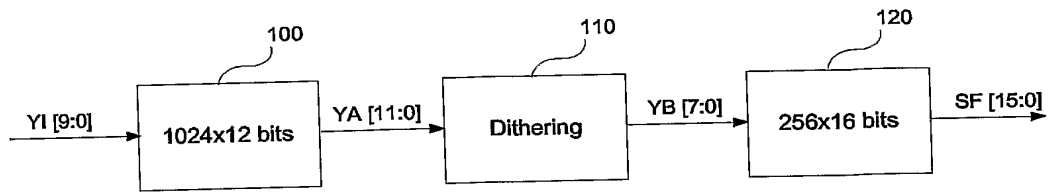


FIG.1 (Prior art)

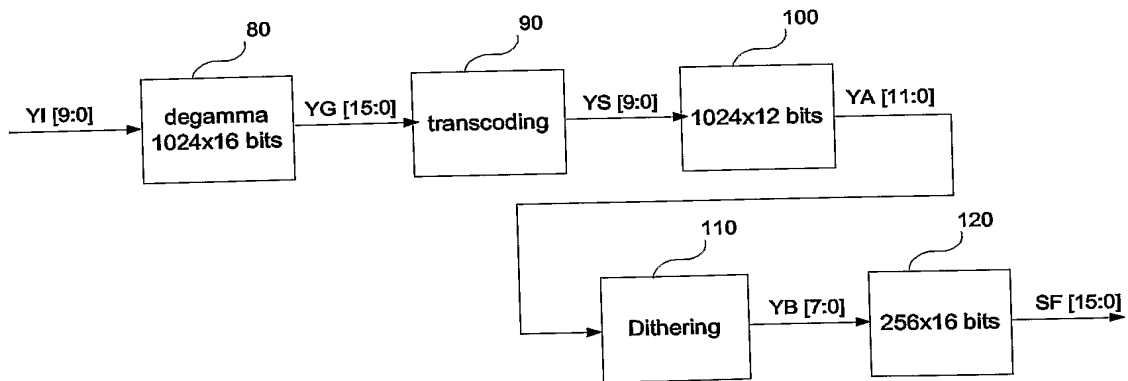


FIG.2

